

# CS170: Discrete Methods in Computer Science

## Summer 2023

### Sets and Friends

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<sup>1</sup>These slides adapt some content from similar slides by Aaron Cote.

# Outline

- 1 Sets
- 2 Tuples and Sequences
- 3 Relations and Functions
- 4 Single-Set Relations

## Definition

A **set** is an unordered collection of distinct objects, which we call it's **members**.

Examples:  $\mathbb{N}$ ,  $\mathbb{Z}$ ,  $\mathbb{R}$ ,  $\mathbb{Q}$ ,  $\emptyset$ , even integers, prime numbers, students in this class, runtime functions that are  $O(n)$

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## Notation

- $\emptyset$  is the empty set
- $\{1, 2, 3\}$ : The set which includes the three number 1, 2, 3
- $\text{Even} = \{x \in \mathbb{N} : \exists k \in \mathbb{N} x = 2k\}$
- $x \in A$  denotes membership. E.g.  $4 \in \text{Even}$
- $x \notin A$  denotes non-membership. E.g.  $3 \notin \text{Even}$

## Note

Order and repetition don't matter!

- E.g.  $\{1, 2, 3\} = \{3, 2, 1\} = \{2, 2, 1, 3, 3, 3\}$

# Relationships between sets

- Subset:  $A \subseteq B$  means every element of  $A$  is in  $B$ 
  - E.g.  $\{1, 2\} \subseteq \{1, 2, 3\}$ ,  $\mathbb{N} \subseteq \mathbb{Z}$ , 170 students  $\subseteq$  USC students
  - The empty set  $\emptyset$  is a subset of every set
  - Every set is a subset of itself: e.g.  $\{1, 2, 3\} \subseteq \{1, 2, 3\}$ .
- Equality:  $A = B$  if both  $A \subseteq B$  and  $B \subseteq A$ .
- Proper subset:  $A \subset B$  or  $A \subsetneq B$  means  $A \subseteq B$  but  $A \neq B$ .
  - E.g. 170 students  $\subset$  USC students,  $\mathbb{N} \subset \mathbb{Z}$
- Superset:  $A \supseteq B$  means  $B \subseteq A$ .
- Proper Superset:  $A \supset B$  or  $A \supsetneq B$  means  $B \subseteq A$  and  $B \neq A$ .
- We say  $A$  and  $B$  are **disjoint** if they have no elements in common
  - E.g. The set of Even numbers and the set of Odd numbers are disjoint

- Sets can include other sets as members. For example
  - $\{\{1\}, \{1, 2, 3\}, \emptyset\}$
  - Set of communities in a social network
  - $\{A \subseteq \mathbb{N} : \sum_{i \in A} i \leq 3\} = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}\}$
  - $\{\emptyset, \{\emptyset\}, \mathbb{N}, \{\mathbb{N}, \mathbb{Q}\}\}$
- **Powerset** of  $A$ , denoted by  $\mathcal{P}(A)$  or  $2^A$ , is set of all subsets of  $A$ 
  - E.g.  $\mathcal{P}(\{1, 2, 3\}) = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$
  - E.g.  $\mathcal{P}(\emptyset) = \{\emptyset\}$ , and  $\mathcal{P}(\{\emptyset\}) = \{\emptyset, \{\emptyset\}\}$

# Cardinality of Sets

- The **cardinality** of a set  $A$ , denoted  $|A|$ , is the number of elements in it. May be finite or infinite. For example:
  - $|\{1, 2, 3\}| = 3$  and  $|\emptyset| = 0$
  - $|\mathcal{P}(\{1, 2, 3\})| = 8$ ,  $|\{\emptyset\}| = 1$ ,  $|\mathcal{P}(\mathcal{P}(\emptyset))| = 2$
  - $|\mathbb{Z}|$  and  $|\mathbb{R}|$  are  $\infty$  (but not the same  $\infty$ !!)
  - $|\{\mathbb{Z}, \mathbb{R}\}| = 2$
- $|\mathcal{P}(A)| = 2^{|A|}$  for finite sets  $A$  (why?)

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## Interesting Fact

Comparing sizes of infinite sets is very interesting and relevant to CS! You will see that cardinality of set of computer programs is a smaller infinity than the cardinality of the set of problems you might want to solve, therefore there are problems that are not computable!



# Operations on Sets and Venn Diagrams

- Intersection:  $A \cap B$  contains elements that are in both  $A$  and  $B$
- Union:  $A \cup B$  contains elements that are in  $A$  or in  $B$  (or both)
- Difference:  $A - B$  or  $A \setminus B$  contains elements that are in  $A$  but not in  $B$
- Complement:  $\overline{A}$  contains elements that are not in  $A$ 
  - Defined relative to a **universe**  $\mathbb{U}$ , which should clear from context.
  - $\overline{A} = \mathbb{U} - A$
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## Some Examples

- $\text{Even} = \mathbb{Z} - \text{Odd}$ , which is  $\overline{\text{Odd}}$  if universe is  $\mathbb{Z}$
- $\text{Even} \cap \text{Odd} = \emptyset$  (they are disjoint)
- $\text{Even} \cup \text{Odd} = \mathbb{Z}$
- $\text{Multiples-of-3} \cap \text{Multiples-of-2} = \text{Multiples-of-6}$
- $\mathbb{Z} \cup \mathbb{R} = \mathbb{R}$
- $\emptyset \cup A = A$  and  $\emptyset \cap A = \emptyset$  for any set  $A$

# Generalized Union and Intersection

Can union or intersect many sets all at once with following shorthand

$$\bigcup_{i=1}^n S_i = S_1 \cup S_2 \cup \dots \cup S_n$$

$$\bigcap_{i=1}^n S_i = S_1 \cap S_2 \cap \dots \cap S_n$$

Can also take infinite union / intersection. For example:

$$\mathbb{N} = \bigcup_{i=1}^{\infty} \{i\}$$

$$A = \bigcap_{x \notin A} (\mathbb{U} - \{x\})$$

# Properties of Set operations

Commutative

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

Distributive

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

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Should remind you of commutative and distributive from propositional logic

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- 3 Relations and Functions
- 4 Single-Set Relations

- An  **$n$ -tuple** is an ordered list of  $n$  elements (basically, an array of  $n$  elements)
  - Written using  $()$  or  $\langle \rangle$ , unlike sets which are written using  $\{\}$
  - E.g.  $(1, 2, 3)$ ,  $(3, 2, 1)$ ,  $(1, 2, 2)$ ,  $(1, 2)$ ,  $(2, 1, 2)$  are all different tuples
  - When  $n = 2$  often called an “ordered pair”
- You often see tuples constructed from sets using **Cartesian products**
- The cartesian product of sets  $A$  and  $B$ , denoted  $A \times B$ , is the set of all ordered pairs  $(a, b)$  with  $a \in A$  and  $b \in B$ .
  - E.g.  $\{1, 2\} \times \{1, 3, 4\} = \{(1, 1), (1, 3), (1, 4), (2, 1), (2, 3), (2, 4)\}$
  - E.g. USC Students  $\times$  USC Courses = Set of possible class enrollments
- Note:  $|A \times B| = |A||B|$

# Sequences

- Intuitively: A sequence is essentially a countably infinite tuple
  - Formally, it is a function from  $\mathbb{N}$  to elements
- E.g. Fibonacci Sequence: 0, 1, 1, 2, 3, 5, 8, ...
- E.g. The sequence  $T(n)$  of worst-case runtimes
- A sequence is called a **recurrence relation** if it is defined recursively
  - E.g. Fibonacci, worst case runtime of Mergesort
- A **closed form expression** for a sequence is an elementary mathematical expression for the  $n$ th element of a sequence
  - We found a closed form expression for the runtime of Mergesort
  - There also is one for the Fibonacci sequence (look it up)
  - Not every sequence has a closed-form expression
  - “closed form” depends on what you allow in your expression



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## Definition

A **Relation** between sets  $A$  and  $B$  is some  $R \subseteq A \times B$ .

## Example

- $A$  is the set of USC Students
- $B$  is the set of USC Courses
- $R = \{(a, b) \in A \times B : \text{student } a \text{ is enrolled in class } b\}$

# Functions

A **function**  $f$  from set  $A$  to set  $B$  takes as input a member of  $a$ , and outputs a member of  $b$ .

## Formally

$f$  is a relation between  $A$  and  $B$  where each  $a \in A$  is related to exactly one  $b \in B$ .

- In other words, each  $a \in A$  shows up exactly once in the relation.
- We use  $f(a)$  to denote the output of  $f$  on  $a$  (i.e., the unique member of  $B$  which is related to  $a$ )
- When  $b = f(a)$ , we say  $b$  is the **image** of  $a$  under  $f$ .
- We say  $f$  is a **map** or **mapping** from  $A$  to  $B$ .
- We call  $A$  the **domain** and  $B$  the **co-domain** of  $f$ .
- The **range** of  $f$  is the set of possible outputs, which may or may not be the entire co-domain
  - $\text{range}(f) = \{f(a) : a \in A\}$ .

# Important Kinds of Functions

A function  $f : A \rightarrow B$  is

- **injective** (a.k.a. **one-to-one**) if different inputs map to different outputs
  - Formally:  $\forall x, y \in A \ (x \neq y \Rightarrow f(x) \neq f(y))$
  - In other words: Every  $b \in B$  is the output of  $f$  on at most one  $a \in A$ .
- **surjective** (a.k.a. **onto**) if every allowed output is produced from some input
  - Formally:  $\forall b \in B \ \exists a \in A \ f(a) = b$
  - In other words: Every  $b \in B$  is the output of  $f$  on at least one  $a \in A$ .
  - In other words still: range = codomain
- **bijective** (a.k.a. one-to-one correspondance) if it is both injective and surjective
  - In other words: Every  $b \in B$  is the output of  $f$  on exactly one  $a \in A$ .

# Examples

- The function mapping USC students to their ID #s is injective, but not surjective onto the co-domain of 10 digit numbers
- $f : \mathbb{R} \rightarrow \mathbb{Z}$  defined by  $f(x) = \lfloor x \rfloor$  is surjective but not injective
- The identity function on  $A$ , defined by  $f(a) = a$ , is bijective for any set  $A$ .
- The function  $f : \mathbb{Z} \rightarrow \text{Even}$  defined by  $f(x) = 2x$  is bijective

## Definition

If  $f : A \rightarrow B$  and  $g : B \rightarrow C$  are functions, then their **composition** is the function  $g \circ f : A \rightarrow C$  is the function mapping  $a \in A$  to  $g(f(a)) \in C$ .

- In other words: Apply  $f$  first, then apply  $g$  to its output.
- **Associative:**  $h \circ (g \circ f) = (h \circ g) \circ f$ 
  - Proof:  $h((g \circ f)(x)) = h(g(f(x))) = (h \circ g)(f(x))$
  - So we omit the parenthesis and write  $h \circ g \circ f$
- **NOT usually commutative:**  $f(g(x)) \neq g(f(x))$ 
  - First of all, this doesn't even type check if  $A \neq C$
  - But even when  $A = B = C$  this isn't true:  $(x + 1)^2 \neq x^2 + 1$

# Inverse of a Function

- If  $f : A \rightarrow B$  is a bijection, then it has an **inverse**  $f^{-1} : B \rightarrow A$ .
  - $f^{-1}(b)$  is defined as the unique  $a \in A$  such that  $f(a) = b$ .
- $f^{-1}(f(a)) = a$  for all  $a \in A$ , and  $f(f^{-1}(b)) = b$  for all  $b \in B$ .
  - **IOW**:  $f^{-1} \circ f$  is the **identity** function on  $A$ , and  $f \circ f^{-1}$  is the identity function on  $B$ .

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  - IOW:  $f^{-1} \circ f$  is the **identity** function on  $A$ , and  $f \circ f^{-1}$  is the identity function on  $B$ .
- When  $f : A \rightarrow B$  is just injective but not surjective, it is common to define  $f^{-1} : \text{range}(f) \rightarrow A$  by first restricting the co-domain to the range (to make it surjective) then taking the inverse of that.



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# Single-Set Relations

Recall:

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- We now look more closely at the special case when  $A = B$ . We call these **Single-set Relations**.
  - A single set relation on  $A$  is some  $R \subseteq A \times A$
- Many relations you have encountered, and will continue to encounter, are on the same set
- We will look at order relations (e.g.  $\leq, <, \subseteq, \subset$ , big- $O$ ) and equivalence relations (e.g.  $=, \equiv$ , big- $\Theta$ )

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## Infix Notation

Often it is convenient to use **infix notation**.

$$a R b \text{ means } (a, b) \in R$$

e.g.  $a \leq b$ ,  $a = b$

# Properties of Single-set Relations

Here are some properties that a single-set relation  $R$  on  $A$  may or may not have:

- **Reflexive**: Every element is related to itself. Formally,  $(x, x) \in R$  for all  $x \in A$ .
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- **Symmetric**: All relationships are mutual. Formally,  $(x, y) \in R \iff (y, x) \in R$  for all  $a, b \in A$ .
- **Antisymmetric**: No relationship between distinct elements is mutual. Formally,  $(x, y) \in R \implies (y, x) \notin R$  for all  $x, y \in A$  with  $x \neq y$ .

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- **Transitive**: If  $x$  is related to  $y$  and  $y$  is related to  $z$  then  $x$  is also related to  $z$ . Formally,  $(x, y) \in R \wedge (y, z) \in R \implies (x, z) \in R$ .

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- **Total**: Every distinct pair of elements is related in at least one direction. Formally,  $(x, y) \in R \vee (y, x) \in R$  for all  $x, y \in A$  with  $x \neq y$ .



# Order Relations

- A **partial order** is a single-set relation which is reflexive, antisymmetric, and transitive
  - E.g.  $\leq$  on numbers,  $\subseteq$  on sets, divisibility  $|$  on integers
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Can be visualized by a directed acyclic graph (DAG)!

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## Topological Ordering

A partial order (weak or strict) can be completed to a total order. For example, the order in which you take your classes, respecting pre-requisites, is a topological ordering of the prerequisite relation.

# Equivalence Relations

## Definition

An **Equivalence Relation** is a single-set relation which is reflexive, symmetric, and transitive.

These capture various notions of “equality”

## Examples

- = on numbers, sets, . . .
- Getting the same grade in 170
- big- $\Theta$  on functions
- $\equiv$  on logical formulas
- Reachability in an undirected graph
- Congruence modulo  $k$ , on integers. Written  $a \equiv b \pmod{k}$ .
  - $a \equiv b \pmod{k}$  iff  $k \mid (a - b)$



# Equivalence Relations

## Useful Fact

An equivalence relation  $\approx$  on a set  $A$  partitions the set into **equivalence classes**, where  $a \approx b$  iff they are in the same class.

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What are the equivalence classes in these examples?

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