# CSCI699: Topics in Learning & Game Theory Lecture 10

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# **Crowdsourcing Information**

In this setting, we have a principal, and n agents. We have a set  $\Omega$  of possible states of nature, where a particular  $\omega \in \Omega$  is drawn from a common prior  $P(\omega)$ . Each agent i receives a signal (a.k.a. type)  $t_i$  from a finite set T of possible types. For any agent i, and any  $t \in T$ ,  $P(t|\omega)$  denotes  $Pr[t_i = t|\omega]$ , thereby,  $t_1, \ldots, t_n$  are conditionally i.i.d. given any  $\omega \in \Omega$ . The goal of the principal is to incentivize agents to report their types truthfully, in order to update his belief about  $\omega$ . That is, he wants to learn the posterior  $P(\omega|t_1,\ldots,t_n)$ .

#### Running Example

Let the state of nature denote whether a new iPhone edition is good, or bad. Accordingly, let  $\Omega = \{G, B\}$  with the associated prior P(G) = 0.6, P(B) = 0.4. Let the types of people denote whether they like or dislike the iPhone, thus, let  $T = \{L, D\}$ . When the iPhone is good, let the type of an agent be distributed as per P(L|G) = 0.75 ( $\Leftrightarrow P(D|G) = 0.25$ ) and, similarly, let the distribution when the iPhone is bad, be P(D|B) = 0.75 ( $\Leftrightarrow P(L|B) = 0.25$ ).

In this case, if n=4 agents report their types to be L,L,L,D respectively, the principal's posterior estimate is  $P(G|L,L,L,D) = \frac{0.6 \cdot 0.75^3 \cdot 0.25}{0.6 \cdot 0.75^3 \cdot 0.25 + 0.4 \cdot 0.25^3 \cdot 0.75} \approx 0.93$  ( $\Leftrightarrow P(B|L,L,L,D) \approx 0.07$ )

This problem can be modeled in several ways.

#### Model 1: Observable Model

In this model, we make the following assumptions.

We assume that the principal solicits the types from agents today,  $\omega$  is directly revealed to the principal tomorrow, and the principal may pay to the agents day after tomorrow.

We also assume that the principal knows  $P(\omega)$  and  $P(t|\omega)$ . (Equivalently, he knows the joint distribution  $P(\omega, t)$ .)

Finally, we make a technical assumption that types  $t \neq t'$  induce distinct posterior distributions w.l.o.g., i.e.  $\exists \omega \in \Omega : P(\omega|t) = P(\omega|t')$ .

With the last assumption, reporting a type t can be considered equivalent to reporting the distribution  $P(\omega|t)$ . Hence, the principal can simply use a strictly proper scoring rule to incentivize the agents to report the respective distributions P(w|t) truthfully. With these truthful reports, with the knowledge of  $P(\omega)$  and  $P(t|\omega)$  as mentioned in the second model assumption, and the inherent property that  $t_i$ 's are i.i.d. given any  $\omega$ , the principal can apply the Bayes' theorem to compute his posterior estimate.

Next, we consider another model which may suit for certain problem scenarios.

### Model 2: Peer Prediction

This model consists of the following assumptions.

As in the case of the Observable model, we assume that the principal knows the joint distribution  $P(\omega, t)$ .

A key difference with the previous model is that the principal never directly sees  $\omega$ . He must decide payments solely based on type reports  $r_1, \ldots, r_n$ .

Finally, we make a technical assumption. Consider  $P(t_j|t_i) = \sum_{\omega} P(\omega|t_i)P(t_j|\omega)$ . We assume that  $P(t_j|t_i=t) \neq P(t_j|t_i=t') \ \forall t \neq t'$ .

To intuitively understand this, consider the iPhone example introduced earlier. One would expect that agent 1 liking iPhone should raise the probability estimate that 2 likes as well. (Note that this holds for the distributions P in general position, even by fixing  $P(t|\omega)$ .)

Now, as per the last assumption, agent i reporting a type is equivalent to him reporting a posterior distribution on the type of some agent  $j \neq i$ . Let  $q_t \in \Delta(T)$  be the conditional distribution of  $t_j$  given  $t_i = t$  for any  $j \neq i$ . (Note that it doesn't depend on i, j.) Then, the *Peer prediction* protocol is as follows:

- 1. To each agent i, assign agent  $\hat{i} \neq i$  as his peer.
- 2. Solicit type reports  $r_1, \ldots, r_n$  from the agents, with each  $r_i \in T$ .

3. Let  $S: \Delta(T) \times T \to \mathbb{R}$  be a strictly proper scoring rule. Pay agent i a value of  $S(q_{r_i}, r_{\hat{i}})$  (computation depends on  $P(\omega)$ ).

For the solution above, we prove the following result.

**Theorem 1.** In the Peer prediction protocol, reporting  $r_i = t_i \ \forall i \ is \ a \ strict \ Bayes-Nash equilibrium.$ 

*Proof.* Assume all agents except i (including  $\hat{i}$ ) report truthfully. Given  $t_i$ , i believes  $t_i \sim q_{t_i}$ ). In step 3 of the protocol, the payment to i is  $S(q_{r_i}, t_{\hat{i}})$  (since we assume  $r_{\hat{i}} = t_{\hat{i}}$ ). Since S is a strictly proper scoring rule, i's best strategy is to report posterior  $q_{t_i}$  on  $t_{\hat{i}}$ , hence he must report  $r_i = t_i$ .

## Model 3: Bayesian Truth Serum

The assumptions here are as follows:

The principal does not know  $P(\omega)$ , nor  $P(t|\omega)$  (even though he knows they exist and that they determine the agents' behavior).

We assume  $n \to \infty$  (This assumption can be removed by using results from follow-up work).

We propose a protocol, called the *Bayesian truth serum* protocol. The idea is as follows. The principal solicits from agent i, not only the type report  $r_i \in T$ , but also a prediction  $y^i \in \Delta(T)$  of the the empirical type frequency  $\bar{x}$  with each  $\bar{x}_j$  defined as the fraction of agents reporting j as their type. Having received  $r_i$ , and  $y^i$  from each agent i, the principal rewards him for a "surprisingly common" type report, i.e., if  $r_i$  is more common in the empirical distribution than estimated by the other agents. Secondly, an agent is also rewarded for a truthful prediction  $y^i$  using a strictly proper scoring rule. Formally,

- 1. Solicit type report  $r_i \in T$ , as well as a prediction  $y^i \in \Delta(T)$  of the the empirical type frequency  $\bar{x}$ .
- 2. Define  $\bar{y}$  so that  $\forall k \in T : \log \bar{y}_k = \frac{1}{n} \sum_i \log y_k^i$ .
- 3. Pay agent i a value of  $F(r_i) + G(y^i)$ , where,  $F(j) = \log \frac{\bar{x}_j}{\bar{y}_j}$  is the reward for "surprisingly common" type reports, and,  $G(y^i) = E_{k \sim \bar{x}}[S(y^i, k)]$  is the reward for accurately predicting  $\bar{x}$  with  $y^i$ , (by choosing S to be a strictly proper scoring rule).

 $G(y^i)$  can be denoted as  $S(y^i, \bar{x})$  as per the notation established in Lecture 9. Clearly, this is maximized at  $y^i = \bar{x}$ .

In the next lecture, we will prove the following result for this protocol.

**Theorem 2.** In the Bayesian truth serum protocol, Each agent i reporting  $r_i = t_i$  and  $y^i$  as their posterior belief on  $\bar{x}$  given  $t_i$ , is a Bayes-Nash equilibrium.