

10/21/10 QIC 710 Lecture 12

Schmidt decomposition

Impossibility of bit commitment

General quantum operations, POVMs, Stinespring dilations

Coin-flipping and bit commitment

- possible based on computational assumptions (one-way permutations)
a guaranteed noisy channel
special relativity [Kent, quant-ph/9906103]
- motivation: rochambeau
secure computation, ...

last time: BB84 bit commitment protocol, not secure

(Alice can change her commitment by keeping a purification of the mixed states sent)

Schmidt decomposition. (for ^{bipartite} pure states only)

Theorem: Let $|\psi\rangle$ be a pure state in $\mathcal{H}_A \otimes \mathcal{H}_B$, $K = \dim \mathcal{H}_A$, $L = \dim \mathcal{H}_B$.

Then there exist orthonormal bases $|e_i\rangle, \dots, |e_K\rangle$ for \mathcal{H}_A
 $|f_i\rangle, \dots, |f_L\rangle$ for \mathcal{H}_B

and reals $\lambda_i \geq 0$ such that

$$|\psi\rangle = \sum_{i=1}^{\min(K,L)} \sqrt{\lambda_i} |e_i\rangle \otimes |f_i\rangle.$$

Proof: $|\psi\rangle = \sum_{i,j} \alpha_{ij} |i\rangle \otimes |j\rangle = \sum_i |i\rangle \otimes |\tilde{i}\rangle$, $|\tilde{i}\rangle = \sum_j \alpha_{ij} |j\rangle$

Choose a basis so that ρ_A is diagonal

$$\rho_A = \sum_{i,j} |i\rangle \langle j| \cdot \text{Tr}(|\tilde{i}\rangle \langle \tilde{j}|)$$

$$\vdots$$

$$\langle \tilde{j} | \tilde{i} \rangle$$

$$\Rightarrow \langle \tilde{j} | \tilde{i} \rangle = 0 \text{ for } i \neq j$$

$$\Rightarrow |\psi\rangle = \sum_i |i\rangle \otimes |\tilde{i}\rangle \quad \checkmark \quad \square$$

$$\Rightarrow \rho_A = \sum_i \lambda_i |e_i\rangle \langle e_i|, \quad \rho_B = \sum_i \lambda_i |f_i\rangle \langle f_i|$$

have the same eigenvalues (nonzero)

Def: $\text{Rank}(\rho_A) = \text{Rank}(\rho_B) :=$ Schmidt number
 a (fragile) entanglement measure (for pure states)

Corollary: Let $|\phi\rangle, |\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B$, $\rho = |\phi\rangle \langle \phi|$, $\sigma = |\psi\rangle \langle \psi|$.

Then $\rho_B = \sigma_B \Leftrightarrow \exists U_A$ so $(U_A \otimes \mathbb{1})|\phi\rangle = |\psi\rangle$.

Proof: Starting with a basis $|f_i\rangle, \dots, |f_L\rangle$ so ρ_B is diagonal,

$$|\phi\rangle = \sum_i \sqrt{\lambda_i} |e_i\rangle \otimes |f_i\rangle$$

$$|\psi\rangle = \sum_i \sqrt{\lambda_i} |e'_i\rangle \otimes |f_i\rangle$$

$$\Rightarrow U_A = \sum_i |e'_i\rangle \langle e_i| \quad \square$$

Corollary: Perfect bit commitment is impossible (think about it).

$$|\phi^0\rangle_{AB} \quad \perp \quad |\phi^1\rangle_{AB}$$

$$\downarrow \quad \quad \downarrow$$

$$\rho_B = \sigma_B \Rightarrow A \text{ is not committed.}$$

Def: A quantum state is a matrix ρ that is
 - positive semi-definite (and Hermitian)
 - with trace 1.

ρ is pure, i.e. $\rho = |\psi\rangle\langle\psi| \Leftrightarrow \text{Tr} \rho^2 = 1$.

unitary evolution: $\rho \mapsto U\rho U^\dagger$

measurement: $\rho \mapsto \frac{\Pi_k \rho \Pi_k}{\text{Tr} \Pi_k \rho}$ with prob. $\text{Tr} \Pi_k \rho$.

General quantum operations: observed dynamics are not always unitary

Def: A general quantum operation (completely positive trace-preserving map) is a mapping

$$\rho \mapsto \sum_k A_k \rho A_k^\dagger$$

where the matrices A_k satisfy

$$\sum_k A_k^\dagger A_k = \mathbb{1}$$

(Kraus operators)

Example 1: Unitary evolution $\rho \mapsto U\rho U^\dagger$

Example 2: Full decoherence (dephasing) of a qubit

$$A_0 = |0\rangle\langle 0|, A_1 = |1\rangle\langle 1|$$

$$\rho = \begin{pmatrix} a & b \\ b^* & 1-a \end{pmatrix} \mapsto \begin{pmatrix} a & 0 \\ 0 & 1-a \end{pmatrix}$$

Example 3: Partial dephasing

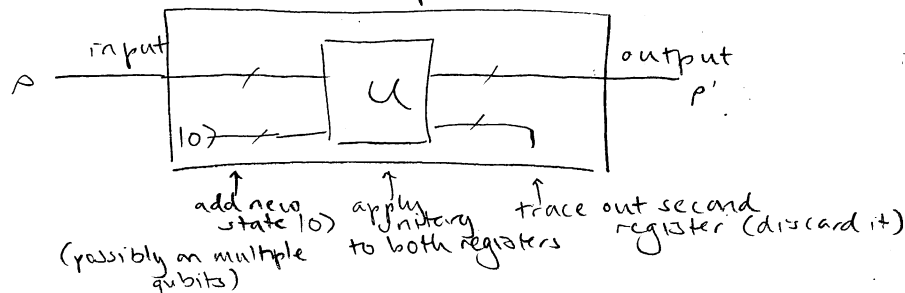
$$A_0 = \sqrt{1-p} \mathbb{1}, A_1 = \sqrt{p} |0\rangle\langle 0|, A_2 = \sqrt{p} |1\rangle\langle 1|$$

$$\rho = \begin{pmatrix} a & b \\ b^* & 1-a \end{pmatrix} \mapsto \begin{pmatrix} a & (1-p)b \\ (1-p)b^* & 1-a \end{pmatrix}$$

time until $p = \frac{1}{e}$ known as T_2 or spin-spin relaxation time
 (more examples later)

Note: The output is a density matrix.

"Derivation": Consider the procedure



We have

$$\begin{aligned}
 \rho' &= \text{Tr}_B \left[U (\rho_A \otimes |0\rangle\langle 0|_B) U^\dagger \right] \\
 &= \sum_k (\mathbb{1} \otimes \langle k|) \left[U (\rho_A \otimes |0\rangle\langle 0|_B) U^\dagger \right] (\mathbb{1} \otimes |k\rangle) \\
 &= \sum_k \underbrace{[(\mathbb{1} \otimes \langle k|) U (\mathbb{1} \otimes |0\rangle)]}_{A_k} \rho \underbrace{[(\mathbb{1} \otimes \langle 0|) U^\dagger (\mathbb{1} \otimes |k\rangle)]}_{A_k^\dagger} \\
 &= \sum_k A_k \rho A_k^\dagger
 \end{aligned}$$

$$\begin{aligned}
 \text{and } \sum_k A_k^\dagger A_k &= \sum_k (\mathbb{1} \otimes \langle 0|) U (\mathbb{1} \otimes |k\rangle) (\mathbb{1} \otimes \langle k|) U^\dagger (\mathbb{1} \otimes |0\rangle) \\
 &= {}_B \langle 0| U \left(\underbrace{\sum_k |k\rangle\langle k|}_\mathbb{1} \right) U^\dagger |0\rangle_B \\
 &= \langle 0| U U^\dagger |0\rangle_B \\
 &= \mathbb{1}_A \quad \checkmark
 \end{aligned}$$

(and in fact only different basis choices do) Observe: Different basis choices give different Kraus reps for same operator. Conversely, any quantum operation can be implemented in this manner. (Stinespring dilation)

Proof:

Given $\{A_k\}_k^m$ Kraus operators satisfying $\sum_k A_k^\dagger A_k = \mathbb{1}_X$. Define U on $\mathcal{H} \otimes \mathbb{C}^m$ by

$$U(|\psi\rangle \otimes |0\rangle) = \sum_k A_k |\psi\rangle \otimes |k\rangle$$

note: Inner products are preserved

$$\begin{aligned}
 \langle \varphi| \otimes \langle 0| U^\dagger U (|\psi\rangle \otimes |0\rangle) \\
 &= \left(\sum_{k'} \langle \varphi| A_{k'}^\dagger \otimes \langle k'| \right) \left(\sum_k A_k |\psi\rangle \otimes |k\rangle \right) \\
 &= \sum_k \langle \varphi| A_k^\dagger A_k |\psi\rangle = \langle \varphi| \psi \rangle
 \end{aligned}$$

so U can be extended to a full unitary.

Moreover,

$$\text{Tr}_B (U |\psi\rangle\langle\psi| \otimes |0\rangle\langle 0| U^\dagger) = \sum_k A_k |\psi\rangle\langle\psi| A_k^\dagger \quad \checkmark \quad \square$$

-allows for simulation of arbitrary quantum operations